

# Macaulay matrices, low displacement rank, and the efficient computation of null spaces

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# Overview

Motivation: Macaulay method for solving (overdetermined) polynomial systems  
Low displacement rank (LDR) matrices and the GKO algorithm  
A fast null space algorithm for Macaulay matrices exploiting LDR properties  
Challenges of extending LDR theory in higher dimensions  
Conclusions and additional remarks

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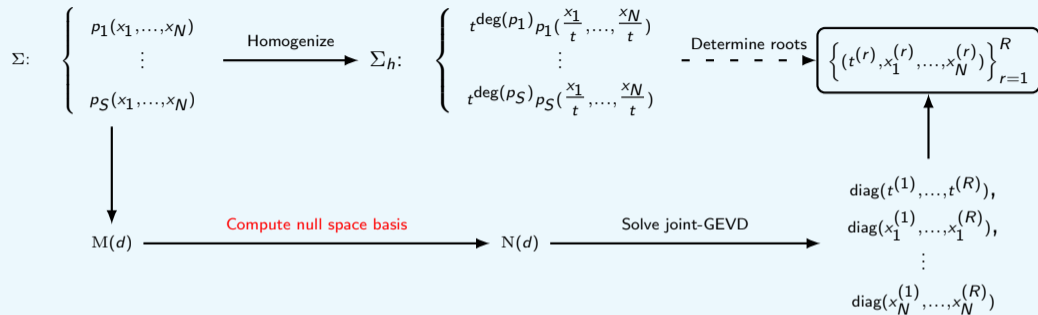
## Solving system of multivariate polynomials

Let  $S \geq N$ , find *isolated* roots of the polynomial system

$$\Sigma : \begin{cases} p_1 = p_1(x_1, x_2, \dots, x_N) \\ \vdots \\ p_S = p_S(x_1, x_2, \dots, x_N) \end{cases} \quad (1)$$

NOTE: In practice, we may be only interested in the affine roots...

## One approach to compute roots (Vanderstukken and De Lathauwer 2021)



The rows of the Macaulay matrix  $M(d)$  span the *set of polynomial combinations*:

$$\left\{ \sum_{s=1}^S h_s \cdot p_s : \deg(h_s) = d - d_s \right\}.$$

## Computing null spaces of Macaulay-type matrices is important!

Null spaces of Macaulay-type matrices are also required in:

1. The GEVD method of *Dreesen, Batselier, and De Moor*.
2. Truncated normal form methods of *Telen, Mourrain, Van Barel*.
3. The Julia package EigenvalueSolver.jl of *Bender and Telen*.
4. *Vermeersch and De Moor's* block-Macaulay method for multiparameter eigenvalue problems.
5. ... and other related methods ...

## Macaulay matrices of univariate systems are block-Toeplitz

$$p_1(x) = 4 + 2x + 1x^2 + 2x^3 + 5x^4$$

$$p_2(x) = 2 + 1x + 4x^2 + 3x^3 + 2x^4$$

$$M(8) = \begin{array}{c} p_1 \\ p_2 \\ xp_1 \\ xp_2 \\ x^2p_1 \\ x^2p_2 \\ x^3p_1 \\ x^3p_2 \\ x^4p_1 \\ x^4p_2 \end{array} \begin{bmatrix} 1 & x & x^2 & x^3 & x^4 & x^5 & x^6 & x^7 & x^8 \\ 4 & 2 & 1 & 2 & 5 & & & & \\ 2 & 1 & 4 & 3 & 2 & & & & \\ & 4 & 2 & 1 & 2 & 5 & & & \\ & 2 & 1 & 4 & 3 & 2 & & & \\ & & 4 & 2 & 1 & 2 & 5 & & \\ & & 2 & 1 & 4 & 3 & 2 & & \\ & & & 4 & 2 & 1 & 2 & 5 & \\ & & & 2 & 1 & 4 & 3 & 2 & \\ & & & & 4 & 2 & 1 & 2 & 5 \\ & & & & 2 & 1 & 4 & 3 & 2 \end{bmatrix}$$

# Macaulay matrices of bivariate systems are (almost) Toeplitz block-(block-)Toeplitz!

$$p_1(x, y) = 1 + 6x + 4x^2 + 2y + 5xy + 3y^2$$

$$p_2(x, y) = 9 + 1x + 3x^2 + 8y + 7xy + 2y^2$$

$$M(4) = \begin{array}{c} \begin{array}{cccccc|cccc|cccc|cc|c} & 1 & x & x^2 & x^3 & x^4 & y & xy & x^2y & x^3y & y^2 & xy^2 & x^2y^2 & y^3 & xy^3 & y^4 \\ p_1 & 1 & 6 & 4 & & & 2 & 5 & & & 3 & & & & & \\ p_2 & 9 & 1 & 3 & & & 8 & 7 & & & 2 & & & & & \\ xp_1 & & 1 & 6 & 4 & & 2 & 5 & & & & 3 & & & & \\ xp_2 & & 9 & 1 & 3 & & 8 & 7 & & & & 2 & & & & \\ x^2p_1 & & & 1 & 6 & 4 & & & 2 & 5 & & & 3 & & & \\ x^2p_2 & & & 9 & 1 & 3 & & & 8 & 7 & & & 2 & & & \\ \hline yp_1 & & & & & & 1 & 6 & 4 & & 2 & 5 & & 3 & & \\ yp_2 & & & & & & 9 & 1 & 3 & & 8 & 7 & & 2 & & \\ xyp_1 & & & & & & & 1 & 6 & 4 & & & 2 & 5 & 3 & \\ xyp_2 & & & & & & & 9 & 1 & 3 & & & 8 & 7 & 2 & \\ \hline y^2p_1 & & & & & & & & & & 1 & 6 & 4 & 2 & 5 & 3 \\ y^2p_2 & & & & & & & & & & 9 & 1 & 3 & 8 & 7 & 2 \end{array} \end{array}$$



## Macaulay matrices of systems with $N \geq 3$ variables

- In trivariate case: Block-Toeplitz where each is Toeplitz-block-(block-)Toeplitz
- For  $N \geq 4$ : higher-order generalizations of the above...

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### LDR matrix of the Sylvester-type

A matrix  $X \in \mathbb{F}^{n \times n}$  has *low displacement rank* with respect to displacement matrices  $A, B \in \mathbb{F}^{n \times n}$  if  $\text{rank } \mathcal{D}_{A,B}\{X\} \ll n$  with

$$\mathcal{D}_{A,B} : X \mapsto AX - XB.$$

One could also work with the Stein equation  $X - AXB$ , leading to different family of algorithms...



## The displacement operator produces a compact representation of the Toeplitz matrix

If  $\varphi \neq 1$ ,  $U$  and  $V$  are generators of the Toeplitz matrix

$$\begin{bmatrix} t_{-3} - t_1 & t_{-2} - t_2 & t_{-1} - t_3 & t_0 - \varphi t_0 \\ 0 & 0 & 0 & t_3 - \varphi t_{-1} \\ 0 & 0 & 0 & t_2 - \varphi t_{-2} \\ 0 & 0 & 0 & t_1 - \varphi t_{-3} \end{bmatrix} = \underbrace{\begin{bmatrix} 1 & 0 \\ 0 & t_3 - \varphi t_{-1} \\ 0 & t_2 - \varphi t_{-2} \\ 0 & t_1 - \varphi t_{-3} \end{bmatrix}}_{=:U} \underbrace{\begin{bmatrix} t_{-3} - t_1 & t_{-2} - t_2 & t_{-1} - t_3 & t_0 - \varphi t_0 \\ 0 & 0 & 0 & 1 \end{bmatrix}}_{=:V^*}$$

Not all generators produce Toeplitz matrices!

LDR matrices associated with  $Z_{n,1}X - XZ_{n,\varphi}$  are called (in this context) Toeplitz-like

As an example,  $T_1^{-1}T_2$  with  $T_1, T_2 \in \mathbb{C}^{n \times n}$  Toeplitz is *Toeplitz-like*.

$$\begin{bmatrix} Z_{n,1} & \\ & Z_{n,1} \end{bmatrix} \begin{bmatrix} T_1 & T_2 \\ -I & 0 \end{bmatrix} - \begin{bmatrix} T_1 & T_2 \\ -I & 0 \end{bmatrix} \begin{bmatrix} Z_{n,\varphi} & \\ & Z_{n,\varphi} \end{bmatrix} \equiv \text{rank } 4$$

↓

$$G_1 \begin{bmatrix} Z_{n,1} & \\ & Z_{n,1} \end{bmatrix} G_1^{-1} G_1 \begin{bmatrix} T_1 & T_2 \\ -I & 0 \end{bmatrix} G_2 - G_1 \begin{bmatrix} T_1 & T_2 \\ -I & 0 \end{bmatrix} G_2 G_2^{-1} \begin{bmatrix} Z_{n,\varphi} & \\ & Z_{n,\varphi} \end{bmatrix} G_2^{-1} \equiv \text{rank } 4$$

↓

$$\begin{bmatrix} Z_{n,1} & \\ * & Z_{n,1} \end{bmatrix} \begin{bmatrix} T_1 & \\ & T_1^{-1}T_2 \end{bmatrix} - \begin{bmatrix} T_1 & \\ & T_1^{-1}T_2 \end{bmatrix} \begin{bmatrix} Z_{n,\varphi} & * \\ & Z_{n,\varphi} \end{bmatrix} \equiv \text{rank } 4$$

↓

$$Z_{n,1} (T_1^{-1}T_2) - (T_1^{-1}T_2) Z_{n,\varphi} \equiv \text{rank } 4$$

## Another example: a pure Cauchy matrix

$$\begin{bmatrix} \mu_1 & & & \\ & \mu_2 & & \\ & & \mu_3 & \\ & & & \mu_4 \end{bmatrix} \begin{bmatrix} \frac{1}{\mu_1 - \nu_1} & \frac{1}{\mu_1 - \nu_2} & \frac{1}{\mu_1 - \nu_3} & \frac{1}{\mu_1 - \nu_4} \\ \frac{1}{\mu_2 - \nu_1} & \frac{1}{\mu_2 - \nu_2} & \frac{1}{\mu_2 - \nu_3} & \frac{1}{\mu_2 - \nu_4} \\ \frac{1}{\mu_3 - \nu_1} & \frac{1}{\mu_3 - \nu_2} & \frac{1}{\mu_3 - \nu_3} & \frac{1}{\mu_3 - \nu_4} \\ \frac{1}{\mu_4 - \nu_1} & \frac{1}{\mu_4 - \nu_2} & \frac{1}{\mu_4 - \nu_3} & \frac{1}{\mu_4 - \nu_4} \end{bmatrix} - \begin{bmatrix} \frac{1}{\mu_1 - \nu_1} & \frac{1}{\mu_1 - \nu_2} & \frac{1}{\mu_1 - \nu_3} & \frac{1}{\mu_1 - \nu_4} \\ \frac{1}{\mu_2 - \nu_1} & \frac{1}{\mu_2 - \nu_2} & \frac{1}{\mu_2 - \nu_3} & \frac{1}{\mu_2 - \nu_4} \\ \frac{1}{\mu_3 - \nu_1} & \frac{1}{\mu_3 - \nu_2} & \frac{1}{\mu_3 - \nu_3} & \frac{1}{\mu_3 - \nu_4} \\ \frac{1}{\mu_4 - \nu_1} & \frac{1}{\mu_4 - \nu_2} & \frac{1}{\mu_4 - \nu_3} & \frac{1}{\mu_4 - \nu_4} \end{bmatrix} \begin{bmatrix} \nu_1 & & & \\ & \nu_2 & & \\ & & \nu_3 & \\ & & & \nu_4 \end{bmatrix} \\
 = \\
 \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$$

If  $\mu_i \neq \mu_j$ , the displacement operator is invertible.

LDR matrices associated with  $\text{diag}(\boldsymbol{\mu})\mathbf{X} - \mathbf{X}\text{diag}(\boldsymbol{\nu})$  are called Cauchy-like

A rank- $r$  *Cauchy-like* matrix  $\mathbf{C} \in \mathbb{C}^{n \times n}$  has entries:

$$[\mathbf{C}]_{ij} = \frac{\mathbf{u}_i^* \mathbf{v}_j}{\mu_i - \nu_j}, \quad \mathbf{u}_i, \mathbf{v}_j \in \mathbb{C}^r.$$

E.g., Loewner matrices with entries  $\frac{\xi_i - \eta_j}{\mu_i - \nu_j}$  are rank-2 *Cauchy-like*.



## The Gohberg-Kailath-Olshevsky (GKO) algorithm for Cauchy-like matrices

$$\begin{bmatrix} \text{diag}(\boldsymbol{\mu}_1) & \\ & \text{diag}(\boldsymbol{\mu}_2) \end{bmatrix} \begin{bmatrix} \mathbf{A} & \mathbf{G}^* \\ \mathbf{F} & \mathbf{B} \end{bmatrix} - \begin{bmatrix} \mathbf{A} & \mathbf{G}^* \\ \mathbf{F} & \mathbf{B} \end{bmatrix} \begin{bmatrix} \text{diag}(\boldsymbol{\nu}_1) & \\ & \text{diag}(\boldsymbol{\nu}_2) \end{bmatrix} = \begin{bmatrix} \mathbf{R}_1 \\ \mathbf{R}_2 \end{bmatrix} \begin{bmatrix} \mathbf{S}_1 \\ \mathbf{S}_2 \end{bmatrix}^*$$

↓

$$\text{diag}(\boldsymbol{\mu}_2) (\mathbf{B} - \mathbf{F}\mathbf{A}^{-1}\mathbf{G}^*) - (\mathbf{B} - \mathbf{F}\mathbf{A}^{-1}\mathbf{G}^*) \text{diag}(\boldsymbol{\nu}_2) = (\mathbf{R}_2 - \mathbf{F}\mathbf{A}^{-1}\mathbf{R}_1) (\mathbf{S}_2 - \mathbf{G}(\mathbf{A}^*)^{-1}\mathbf{S}_1)^*$$

### Main principles behind $O(n^2)$ GKO algorithm:

- *Schur complements* preserve Cauchy → use *generators* for Gauss elimination.
- *Permutations* preserve Cauchy → partial pivoting for improved numerical stability

Stability of GKO can be further enhanced through *approximate* complete pivoting

$$\text{diag}(\boldsymbol{\mu})\mathbf{C} - \mathbf{C}\text{diag}(\boldsymbol{\nu}) =: \mathbf{G} = \begin{bmatrix} | & & | \\ \mathbf{q}_1 & \cdots & \mathbf{q}_r \\ | & & | \end{bmatrix} \begin{bmatrix} w_{11} & w_{12} & \cdots & w_{1n} \\ \vdots & \vdots & & \vdots \\ w_{11} & w_{12} & \cdots & w_{1n} \end{bmatrix}.$$

(Gu 1998, Lemma 3.1)

The column with the *largest* 2-norm in  $\mathbf{G}$  (denoted  $j_{\max}$ ) satisfies:

$$\max_{1 \leq i \leq n} |c_{ij_{\max}}| \geq \frac{1}{K(\boldsymbol{\mu}, \boldsymbol{\nu})\sqrt{n}} \max_{1 \leq i, j \leq n} |c_{ij}|.$$

If  $r \ll n$  and  $\mathbf{Q}$  is orthonormal, finding  $j_{\max}$  is cheap

Gauss updates *destroy* orthonormality  $\rightarrow$  **extra cost:** QR-decomposition at each step!

Many LDR matrices are efficiently converted to Cauchy-like (Heinig 1995)

Assume  $A, B$  have "fast" eigendecompositions

$$AX - XB = UV^*$$

↓

$$(W_A \Lambda_A W_A^{-1}) X - X (W_B \Lambda_B W_B^{-1}) = UV^*$$

↓

$$\Lambda_A W_A^{-1} X W_B - W_A^{-1} X W_B \Lambda_B = (W_A^{-1} U) (W_B^* V)^*$$

↓

$$\text{diag}(\mu) C - C \text{diag}(\nu) = R S^*$$

$Z_{n,\varphi}$  has a fast eigendecomposition, thus Toeplitz-like systems are solved fast using GKO!

Denote  $\omega_n := \exp(-2\pi i/n)$  and  $F_n \in \mathbb{C}^{n \times n}$  the (unitary) DFT matrix. Then

$$Z_{n,\varphi} = (D_{n,\varphi} F_n) (\varphi^{1/n} \Omega_n) (D_{n,\varphi} F_n)^{-1},$$

where  $D_{n,\varphi} := \text{diag}(1, \varphi^{-1/n}, \dots, \varphi^{-(n-1)/n})$ ,  $\Omega_n := \text{diag}(1, \bar{\omega}_n, \dots, \bar{\omega}_n^{n-1})$ .

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## $M(d)$ is LDR for univariate systems

Assuming  $\deg(p_s) = d_\Sigma$  for  $s = 1, 2, \dots, S$ ,

$$M(d) := \begin{bmatrix} \mathbf{c}_{0i} & \mathbf{c}_{1i} & \cdots & \mathbf{c}_{(d_\Sigma-i)i} & & & & \\ & \mathbf{c}_{0i} & \mathbf{c}_{1i} & \cdots & \mathbf{c}_{(d_\Sigma-i)i} & & & \\ & & \ddots & \ddots & & \ddots & & \\ & & & \mathbf{c}_{0i} & \mathbf{c}_{1i} & \cdots & \mathbf{c}_{(d_\Sigma-i)i} & \end{bmatrix} \in \mathbb{C}^{S(\Delta d+1) \times (d+1)}.$$

For  $\mathcal{D} : X \mapsto (Z_{d-d_\Sigma,1} \otimes I_S) X - XZ_{d+1,\varphi}$ , we have

$$\text{rank } \mathcal{D} \{M(d)\} = 2S, \quad \text{irrespective of } d.$$

## The Macaulay matrix for the general bivariate case

Assuming  $\deg(p_s) = d_\Sigma$  for  $s = 1, 2, \dots, S$ ,

$$M(d) := \begin{bmatrix} M_{0,0} & M_{1,0} & \cdots & M_{d_\Sigma,0} \\ & M_{0,1} & M_{1,1} & \cdots & M_{d_\Sigma,1} \\ & & \ddots & \ddots & \ddots \\ & & & M_{0,\Delta d} & M_{1,\Delta d} & \cdots & M_{d_\Sigma,\Delta d} \end{bmatrix} \in \mathbb{C}^{\frac{S}{2}(\Delta d+1)(\Delta d+2) \times \frac{1}{2}(d+1)(d+2)},$$

with  $\Delta d := d - d_\Sigma$  and

$$M_{i,j} := \begin{bmatrix} \mathbf{c}_{0i} & \mathbf{c}_{1i} & \cdots & \mathbf{c}_{(d_\Sigma-i)i} \\ & \mathbf{c}_{0i} & \mathbf{c}_{1i} & \cdots & \mathbf{c}_{(d_\Sigma-i)i} \\ & & \ddots & \ddots & \ddots \\ & & & \mathbf{c}_{0i} & \mathbf{c}_{1i} & \cdots & \mathbf{c}_{(d_\Sigma-i)i} \end{bmatrix} \in \mathbb{C}^{S(\Delta d+1-j) \times (d+1-i-j)}.$$

$M(d)$  is *relatively* LDR for bivariate systems

$$\mathcal{D} \{M(d)\} = \begin{bmatrix} Z_{\Delta d+1,1} \otimes I_S & & \\ & \ddots & \\ & & Z_{1,1} \otimes I_S \end{bmatrix} M(d) - M(d) \begin{bmatrix} Z_{d+1,\varphi_{d+1}} & & \\ & \ddots & \\ & & Z_{1,\varphi_1} \end{bmatrix}$$

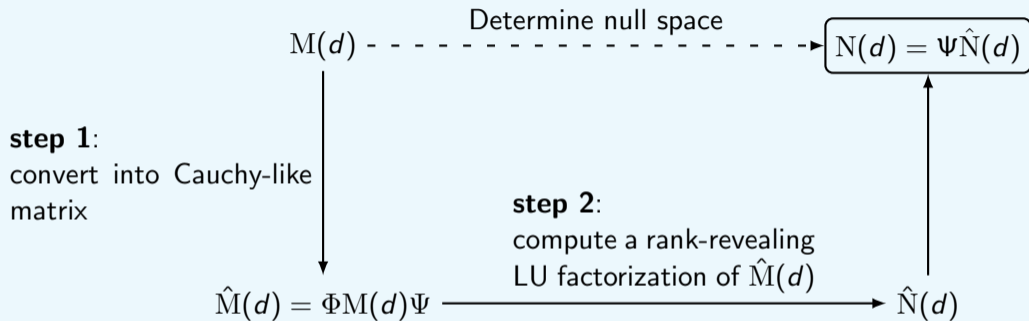
Dimensions of  $M(d) \in \mathbb{C}^{\frac{S}{2}(\Delta d+1)(\Delta d+2) \times \frac{1}{2}(d+1)(d+2)}$  grow *quadratically* w.r.t.  $d$ , but

$$\text{rank } \mathcal{D} \{M(d)\} \leq S(\Delta d + 1) = S(d + 1 - d_\Sigma).$$

grows only *linearly* with  $d$ .



## An overview of the fast algorithm for computing the null space



Both steps can be done *fast*!

## Rank-revealing LU (RRLU) factorization (Miranian and Gu 2003)

Apply permutation matrices  $\Pi_1, \Pi_2$  such that  $\hat{M}_{22} - \hat{M}_{21}\hat{M}_{11}^{-1}\hat{M}_{12} \approx 0$  in

$$\begin{aligned}\Pi_1 \hat{M}(d) \Pi_2 &= \begin{bmatrix} \mathbf{I} & \\ \hat{M}_{21}\hat{M}_{11}^{-1} & \mathbf{I} \end{bmatrix} \begin{bmatrix} \hat{M}_{11} & \\ & \hat{M}_{22} - \hat{M}_{21}\hat{M}_{11}^{-1}\hat{M}_{12} \end{bmatrix} \begin{bmatrix} \mathbf{I} & \hat{M}_{11}^{-1}\hat{M}_{12} \\ & \mathbf{I} \end{bmatrix} \\ &\approx \begin{bmatrix} \hat{M}_{11} \\ \hat{M}_{21} \end{bmatrix} \begin{bmatrix} \mathbf{I} & \hat{M}_{11}^{-1}\hat{M}_{12} \end{bmatrix}\end{aligned}$$

Expression for null space:

$$\hat{N}(d) = \Pi_2 \begin{bmatrix} \tilde{N} \\ \mathbf{I} \end{bmatrix}, \quad \tilde{N} := -\hat{M}_{11}^{-1}\hat{M}_{12}.$$

## Extending the GKO algorithm to compute null spaces

- $\tilde{N}$  is also Cauchy-like of the same LDR rank as  $\hat{M}(d)$
- Adapt GKO algorithm to directly determine generators for  $\tilde{N}$
- Use (approximate) complete pivoting strategy!
- Bivariate systems: QR too expensive  $\rightarrow$  update strategy using Householder!
- Special provisions for stable calculation of  $\tilde{N}$ .
- For more technical details (Govindarajan, Widdershoven, et al. 2024)

Flop count is reduced from  $\mathcal{O}(d_\Sigma^6)$  to  $\mathcal{O}(d_\Sigma^5)$  for bivariate systems

Assumptions:

- Number of roots  $\sim d_\Sigma^2$  (i.e., Bezout bound),
- $2 \leq S \ll d_\Sigma$ .

A quick complexity overview for each step

- Step 1:  $\mathcal{O}(S \cdot d_\Sigma \cdot \Delta d \cdot d \log d)$
- Step 2:  $\mathcal{O}((\text{rank } M(d)) \cdot S^2 d^3)$ ,  $\text{rank } M(d) \sim d^2$

$d \leq 2d_\Sigma - 2$  to find a null space containing all system roots

$$\downarrow$$
$$\mathcal{O}(d_\Sigma^5)$$

## Stability experiments: error grows linearly with problem size

$$\epsilon := \frac{\|M(d)Q\|_2}{\|M(d)\|_2}, \quad Q \text{ is an orthonormal basis for } \text{col } N(d)$$

	$d_\Sigma$				
	2	4	8	16	32
SVD on $M(d)$	2.23e-16	3.75e-16	5.70e-16	7.94e-16	9.51e-16
SVD on $\hat{M}(d)$	2.57e-16	4.77e-16	7.54e-16	9.97e-16	1.15e-15
GECP on $M(d)$	1.40e-16	3.11e-16	8.33e-16	1.02e-14	1.40e-13
GECP on $\hat{M}(d)$	2.08e-16	4.65e-16	1.03e-15	9.73e-15	1.21e-13
GECP on $\mathcal{C}$	4.35e-16	1.51e-15	1.35e-14	1.72e-13	2.81e-12
GEAP on $\mathcal{C}$	4.21e-16	3.63e-15	3.88e-14	3.19e-13	4.48e-12

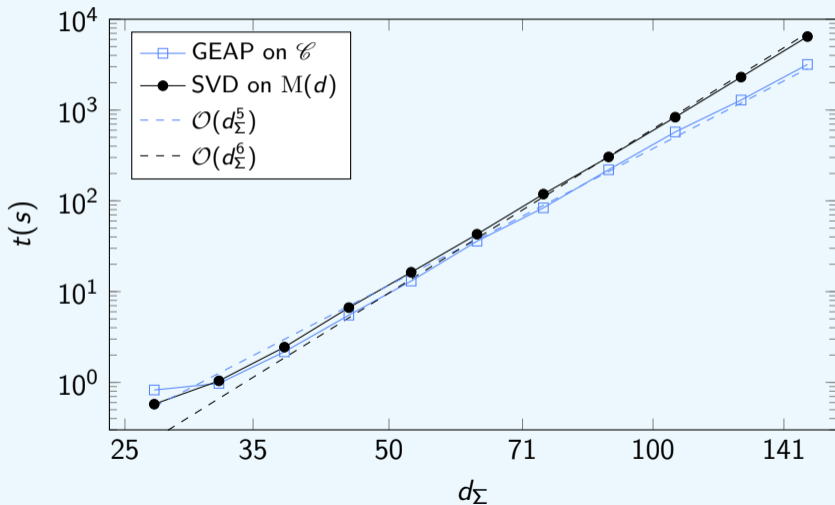
Median error  $\epsilon$  over 100 runs for *square* systems with different methods and degrees.

### Sources of error:

- switching to LU instead of an SVD
- working with the compact Cauchy representation  $\mathcal{C}$
- switching to approximate pivoting ← Surprisingly not so bad!

Our experiments indicate that the flop complexity is indeed  $O(d_\Sigma^5)$

The measurements are the median of an adapted number of runs after warmup.



The fast algorithm extends to Polynomial systems in the Chebyshev basis

$$\Sigma : \left\{ \begin{array}{l} p_1(x, y) := \sum_{i=0}^{d_\Sigma} \sum_{j=0}^{d_\Sigma-i} b_{1ij} T_i(x) T_j(y) = 0 \\ \vdots \\ p_S(x, y) := \sum_{i=0}^{d_\Sigma} \sum_{j=0}^{d_\Sigma-i} b_{Sij} T_i(x) T_j(y) = 0 \end{array} \right.$$

- *Toeplitz-plus-Hankel* structures instead of just Toeplitz.
- Apply same techniques but with a *modified* displacement equation.

## One-variable case

Since  $T_k(x)T_l(x) = \frac{1}{2}(T_{k+l}(x) + T_{|k-l|}(x))$ , we have:

$$W(d) = W^{\text{tpz}}(d) + W^{\text{hnk}}(d) \in \mathbb{C}^{S(\Delta d+1) \times d}$$

$$W^{\text{tpz}}(d) := \begin{bmatrix} \mathbf{b}_{0j} & \mathbf{b}_{1j} & \cdots & \mathbf{b}_{d_{\Sigma}j} \\ \mathbf{b}_{1j} & \mathbf{b}_{0j} & \mathbf{b}_{1j} & \cdots & \mathbf{b}_{d_{\Sigma}j} \\ \vdots & \ddots & \ddots & \ddots & \ddots \\ \mathbf{b}_{d_{\Sigma}j} & \cdots & \mathbf{b}_{1j} & \mathbf{b}_{0j} & \mathbf{b}_{1j} & \cdots & \mathbf{b}_{d_{\Sigma}j} \\ & \ddots & & \ddots & \ddots & \ddots & \ddots \\ & & \mathbf{b}_{d_{\Sigma}j} & \cdots & \mathbf{b}_{1j} & \mathbf{b}_{0j} & \mathbf{b}_{1j} & \cdots & \mathbf{b}_{d_{\Sigma}j} \end{bmatrix}, \quad W^{\text{hnk}}(d) := \begin{bmatrix} \mathbf{b}_{0j} & \mathbf{b}_{1j} & \cdots & \mathbf{b}_{d_{\Sigma}j} \\ \mathbf{b}_{1j} & & \ddots & \\ \vdots & & \ddots & \\ \mathbf{b}_{d_{\Sigma}j} & & & \end{bmatrix}$$

$$\text{rank}((Y_{\Delta d,0} \otimes I_S)W(d) - W(d)Y_{d,1}) = 4S, \quad Y_{n,\delta} := \begin{bmatrix} \delta & 1 & & \\ 1 & 0 & \ddots & \\ & 1 & \ddots & 1 \\ & & \ddots & 0 & 1 \\ & & & 1 & \delta \end{bmatrix}$$



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The algorithm does not nicely extend for general  $N$ -variable systems

- The idea of “GKO+complete pivoting” naturally extends to  $N$ -variable system
- Current set-up however has diminishing returns:  $\mathcal{O}(d^{3N})$  to  $\mathcal{O}(d^{3N-1})$ .
- Issue: we only exploit one direction of the multi-Toeplitz structure!

How to exploit multi-level Toeplitz structures? A big open question!

## What we are searching for!

Given a Toeplitz-block toeplitz matrix  $T \in \mathbb{C}^{n^2 \times n^2}$ , does there exist an  $A, B \in \mathbb{C}^{n^2 \times n^2}$  such that

$$\text{rank}(AT - TB) \in O(1)$$

and  $A, B$  have “fast” eigendecompositions?

*If answer is no, we need a radically different approach...*

# Could composing displacement operators possibly resolve things?

$$T = \begin{bmatrix} \begin{matrix} \sigma_3 \\ \sigma_2 \\ \sigma_1 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \\ \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \\ \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \\ \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} & \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \begin{matrix} \sigma_1 \\ \sigma_2 \\ \sigma_3 \end{matrix} \end{bmatrix}$$

$$D_1 : T \mapsto (I_n \otimes Z_{n,\varphi_a}) T - T (I_n \otimes Z_{n,\varphi_b})$$

$$D_2 \circ D_1 : D_1\{T\} \mapsto (Z_{n,\varphi_c} \otimes I_n) D_1\{T\} - D_1\{T\} (Z_{n,\varphi_d} \otimes I_n)$$

## Could composing displacement operators possibly resolve things?

$$\mathcal{D}_1\{\mathbf{T}\} = \begin{bmatrix} * & * & * & * & | & * & * & * & * & | & * & * & * & * & | & * & * & * & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ * & * & * & * & | & * & * & * & * & | & * & * & * & * & | & * & * & * & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ * & * & * & * & | & * & * & * & * & | & * & * & * & * & | & * & * & * & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ * & * & * & * & | & * & * & * & * & | & * & * & * & * & | & * & * & * & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \\ 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * & | & 0 & 0 & 0 & * \end{bmatrix} \sim \text{rank} - O(n) \text{ matrix}$$

$$\mathcal{D}_1 : \quad \mathbf{T} \mapsto (\mathbf{I}_n \otimes \mathbf{Z}_{n,\varphi_a}) \mathbf{T} - \mathbf{T} (\mathbf{I}_n \otimes \mathbf{Z}_{n,\varphi_b})$$

$$\mathcal{D}_2 \circ \mathcal{D}_1 : \quad \mathcal{D}_1\{\mathbf{T}\} \mapsto (\mathbf{Z}_{n,\varphi_c} \otimes \mathbf{I}_n) \mathcal{D}_1\{\mathbf{T}\} - \mathcal{D}_1\{\mathbf{T}\} (\mathbf{Z}_{n,\varphi_d} \otimes \mathbf{I}_n)$$



Generally, composition displacement operators yields generalized Sylvester equations

If  $\mathcal{D}_{A,B} : T \mapsto AT - TB$  and  $\mathcal{D}_{C,D} : T \mapsto CT - TD$ , then

$$\mathcal{D}_{C,D} \circ \mathcal{D}_{A,B} : T \mapsto CAT - CTB - ATD + TBD$$

## Solving generalized Sylvester equations

Solve

$$\sum_{i=1}^p (A_i \otimes B_i) \text{vec}(X) = \text{vec}(Y) \iff \sum_{i=1}^p A_i X B_i = Y$$

Diagonalization of the system:

- For  $p = 2$ , substitute gen. eigendecomp. of pencils  $(A_1, A_2)$  and  $(B_1, B_2)$ :

$$\begin{aligned} A_1 X B_1 + A_2 X B_2 &= Y \\ V_A \Lambda_{A_1} Q_A^{-1} X V_B \Lambda_{B_1} Q_B^{-1} + V_A \Lambda_{A_2} Q_A^{-1} X V_B \Lambda_{B_2} Q_B^{-1} &= Y \\ \Lambda_{A_1} Q_A^{-1} X V_B \Lambda_{B_1} + \Lambda_{A_2} Q_A^{-1} X V_B \Lambda_{B_2} &= V_A^{-1} Y Q_B \end{aligned}$$

- For  $p > 2$  we require  $\{A_i\}_{i=1}^p$  and  $\{B_i\}_{i=1}^p$  to be both jointly diagonalizable.



Could this possibly be true?

Given a Toeplitz-block-Toeplitz matrix  $T \in \mathbb{C}^{n^2 \times n^2}$ , does there exist for some  $p > 2$   $\{A_i\}_{i=1}^p$  and  $\{B_i\}_{i=1}^p$  such that

$$\text{rank} \left( \sum_{i=1}^p A_i T B_i \right) \in O(1)$$

and  $\{A_i\}_{i=1}^p, \{B_i\}_{i=1}^p$  are jointly diagonalizable by “nice” matrices?

*If answer is no, we again need a radically different approach...*

# Overview

Motivation: Macaulay method for solving (overdetermined) polynomial systems

Low displacement rank (LDR) matrices and the GKO algorithm

A fast null space algorithm for Macaulay matrices exploiting LDR properties

Challenges of extending LDR theory in higher dimensions

**Conclusions and additional remarks**

## Conclusions and additional remarks






The Macaulay null space problem:

- The GKO algorithm allows for a *faster* algorithm than standard approaches.
- But the returns *diminish* as  $N$  becomes larger.
- Other LDR methods (Mastronardi, Van Barel, et al. 2009) suffer from *same* issue.
- Further improvements possible with HSS representations of Cauchy-like matrices
- They sadly come with large constants and do *not* resolve dimensionality issues.
- Root cause: *inability* to fully exploit multi-Toeplitz structure in Gauss elimination
- Any breakthrough on this front will have *great impact*, e.g., faster TBT solvers!



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